

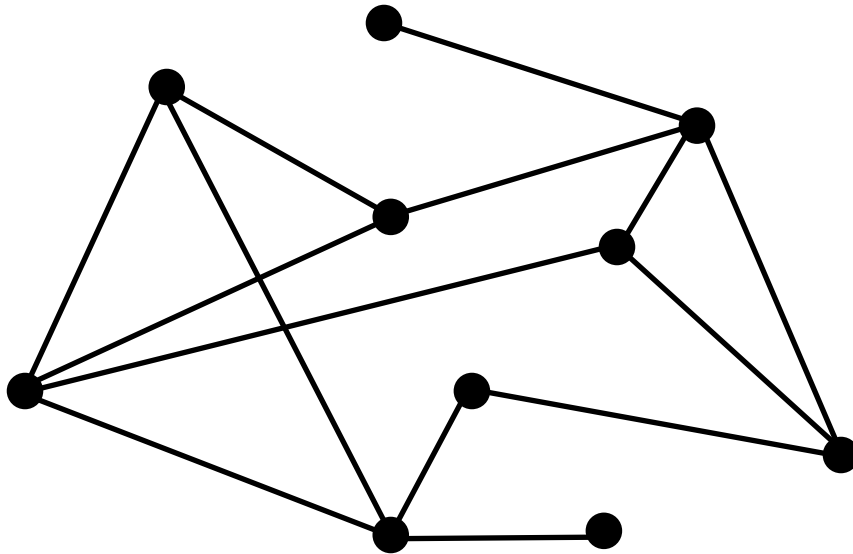
Coloring intersection points of line segments

Sam Lowery

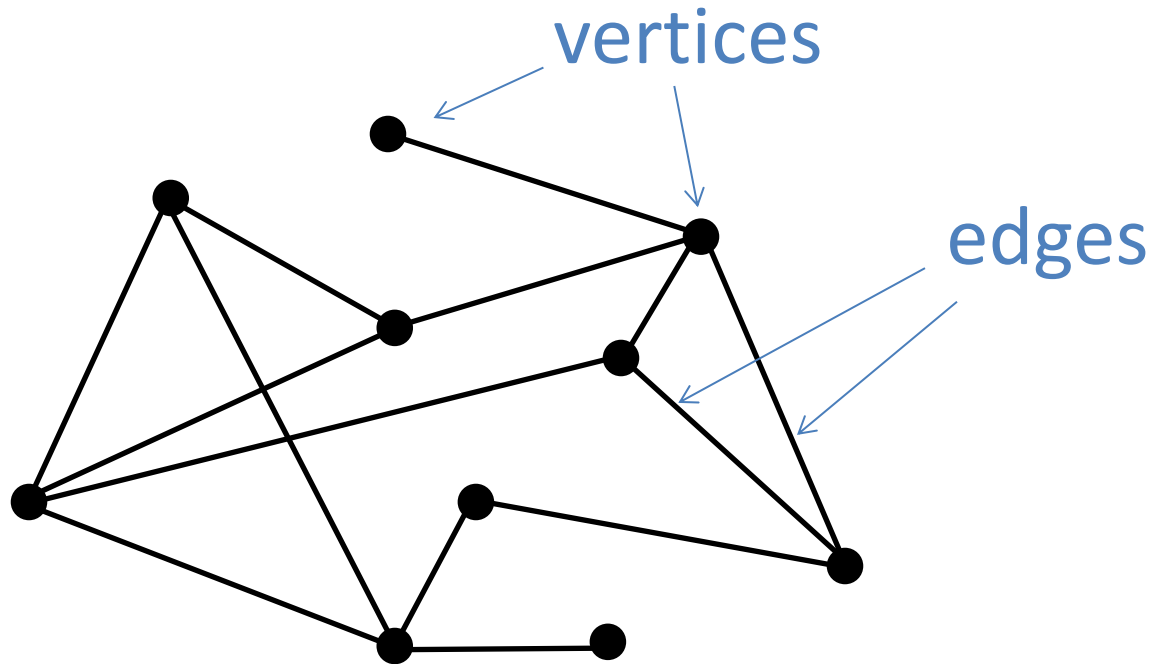
Department of Mathematics & Statistics

2022 Symposium for Student Research,
Scholarship, and Creative Activity

Graphs



Graphs



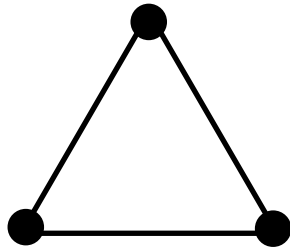
Complete graphs



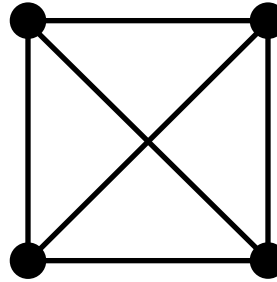
K_1



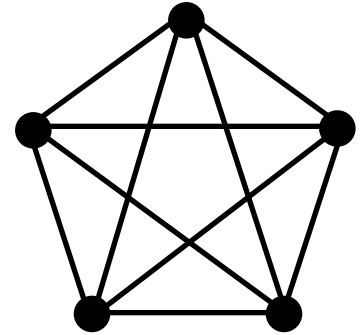
K_2



K_3



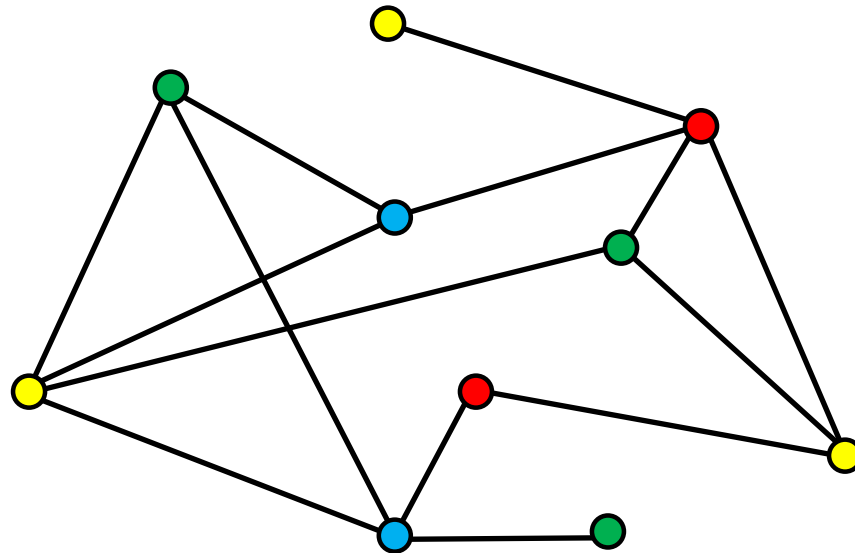
K_4



K_5

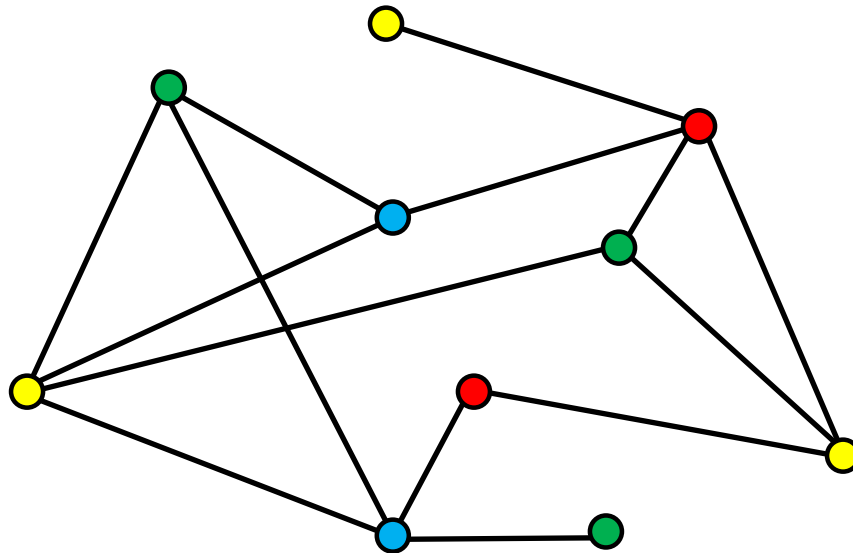
Coloring a graph

Coloring a graph means coloring its vertices so that adjacent vertices get different colors



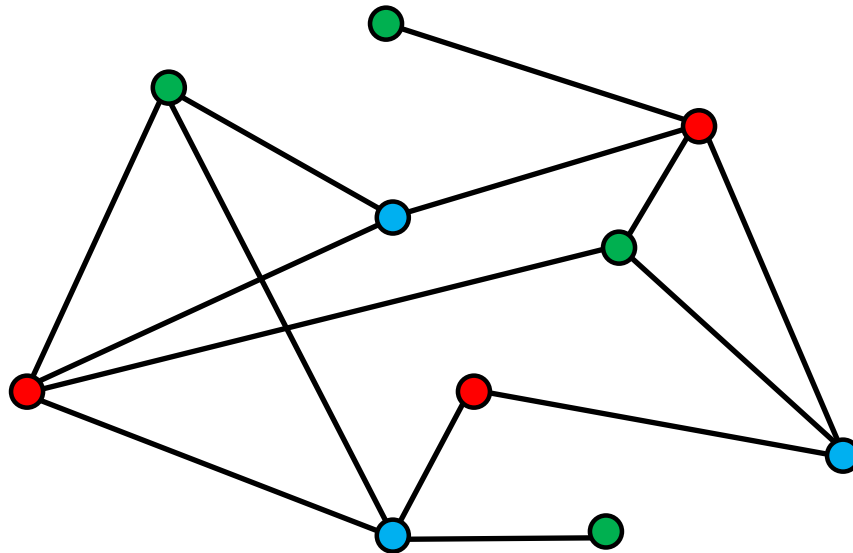
Coloring a graph

$\chi(G)$: smallest possible number of colors with which a graph G can be colored



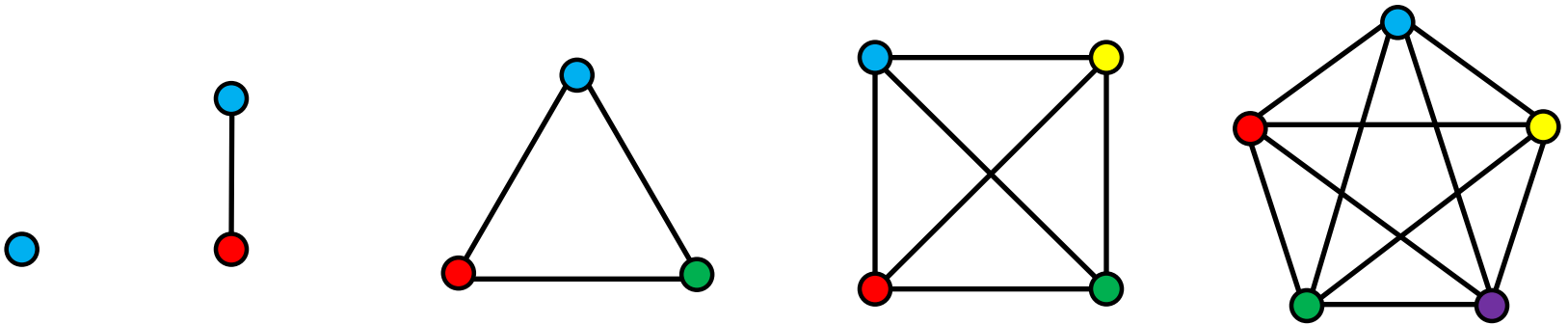
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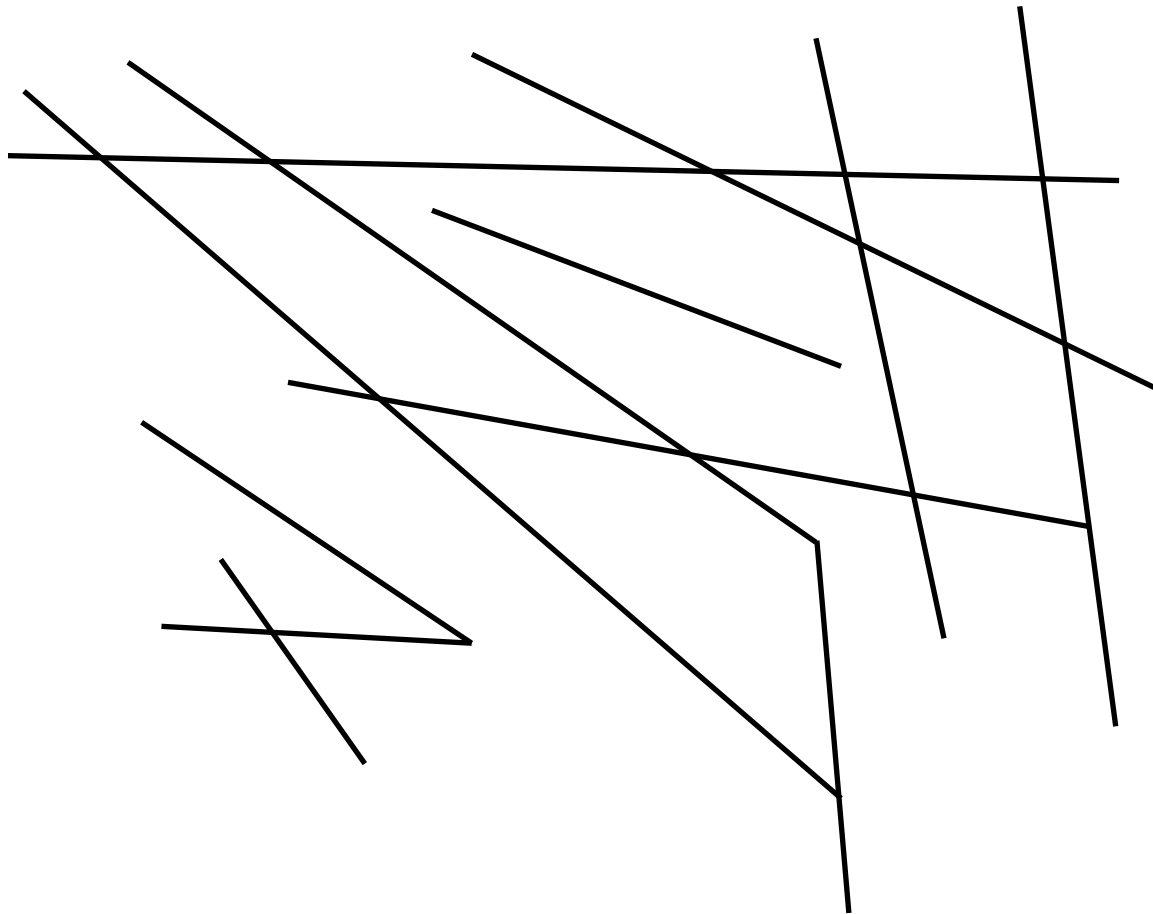
$$\chi(G) = 3$$

Coloring a graph



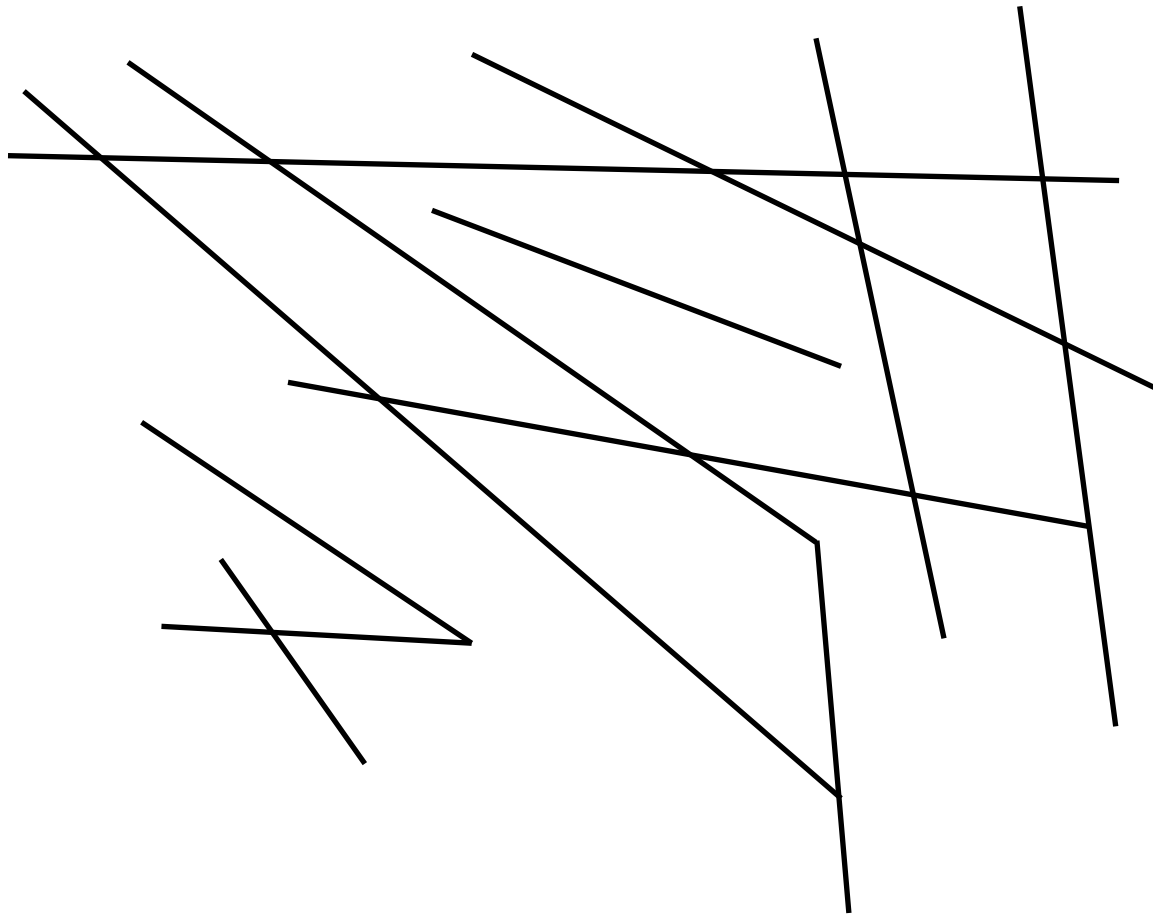
$$\chi(K_n) = n$$

Set of line segments



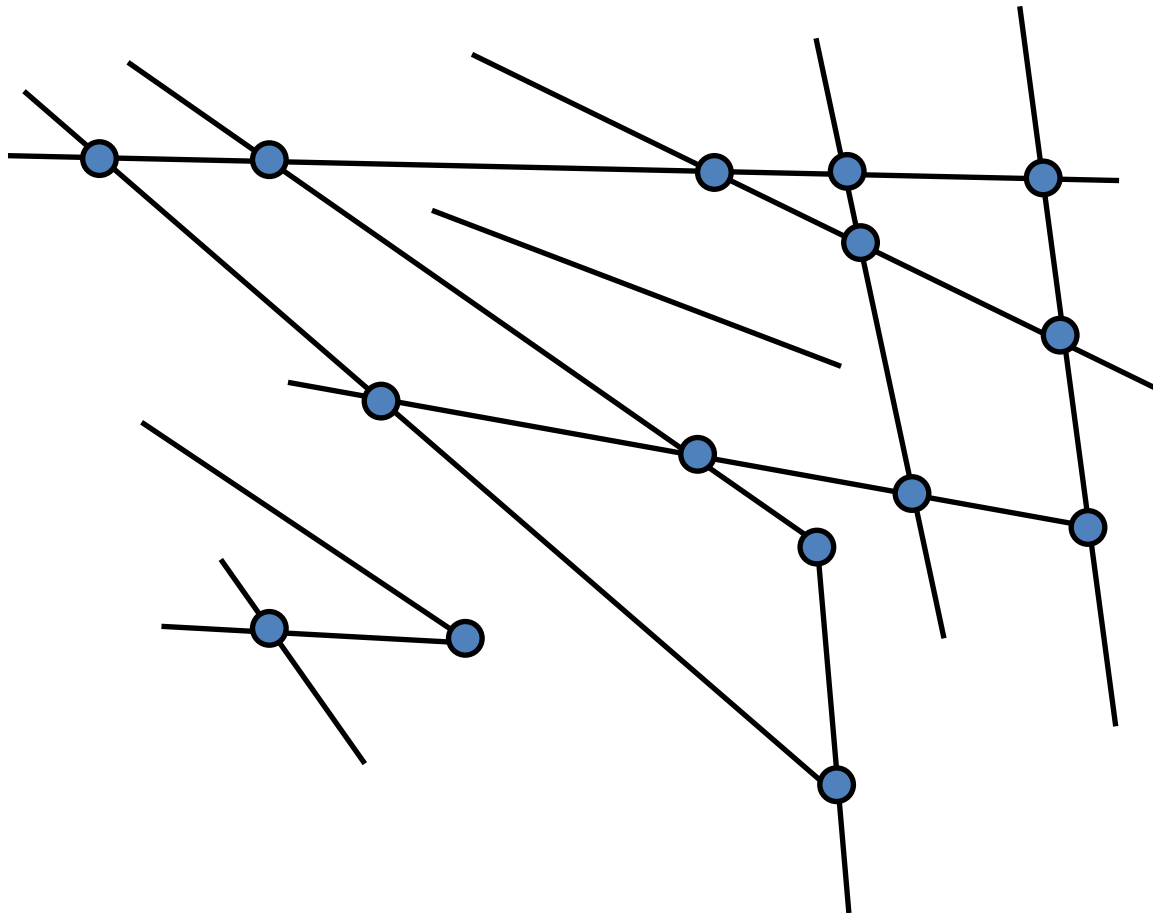
Set of line segments

M



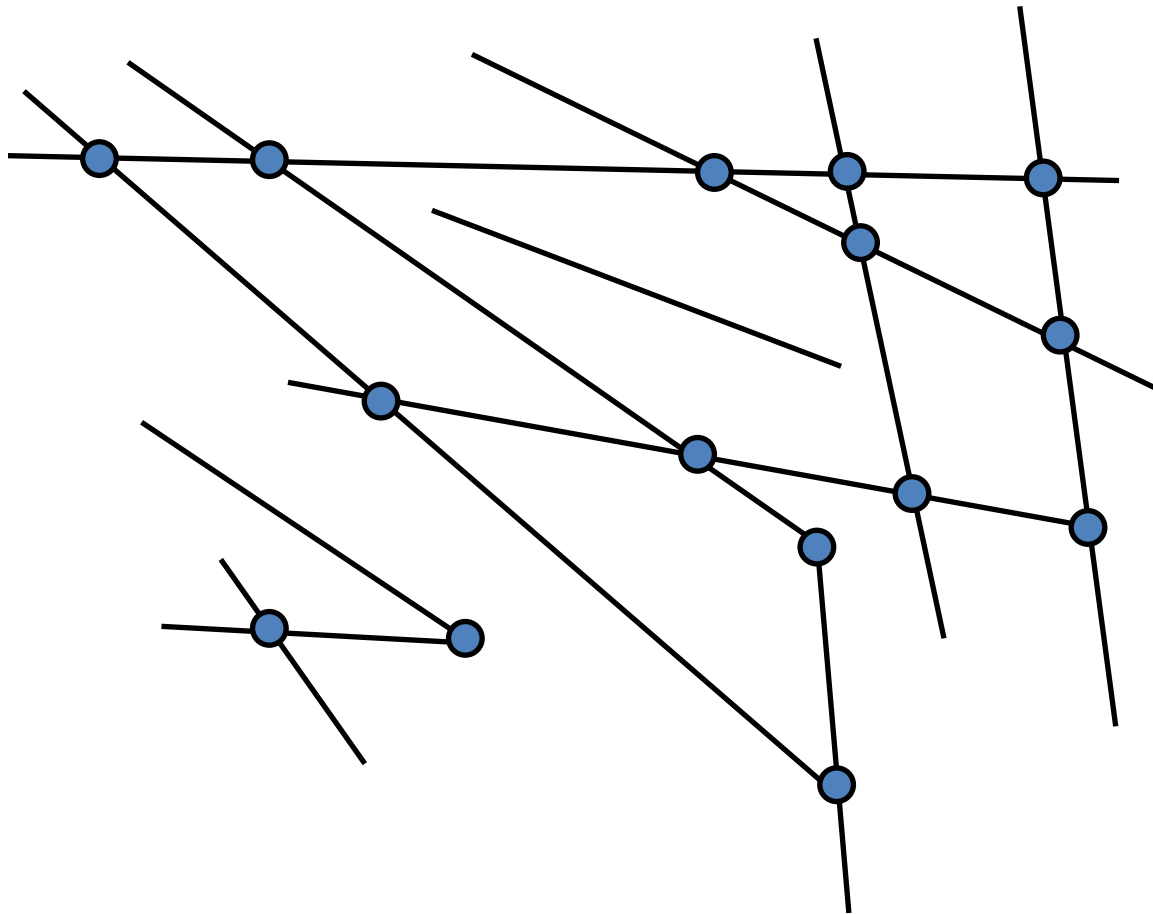
Intersection points

$P(M)$



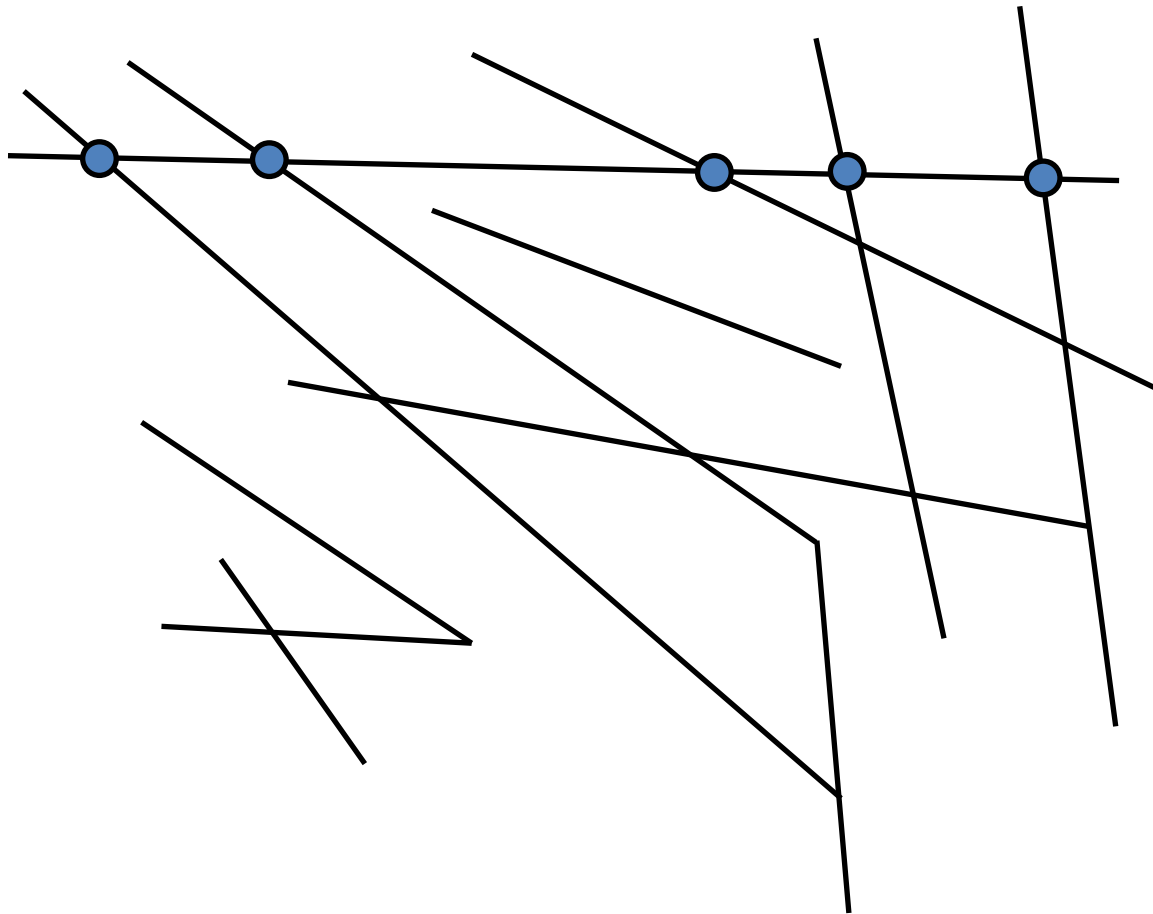
Intersection points

$w(M)$: most intersection points in a segment



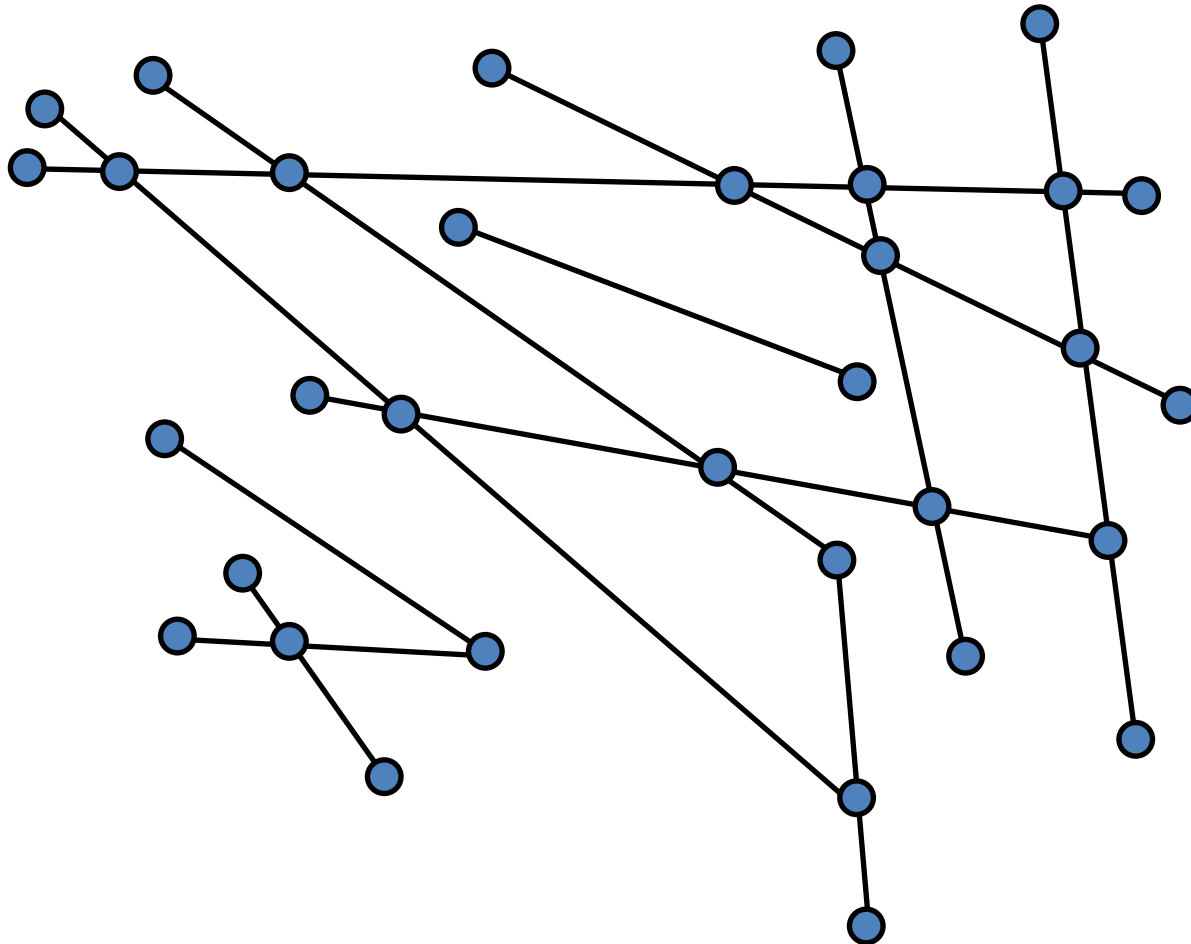
Intersection points

$$w(M) = 5$$



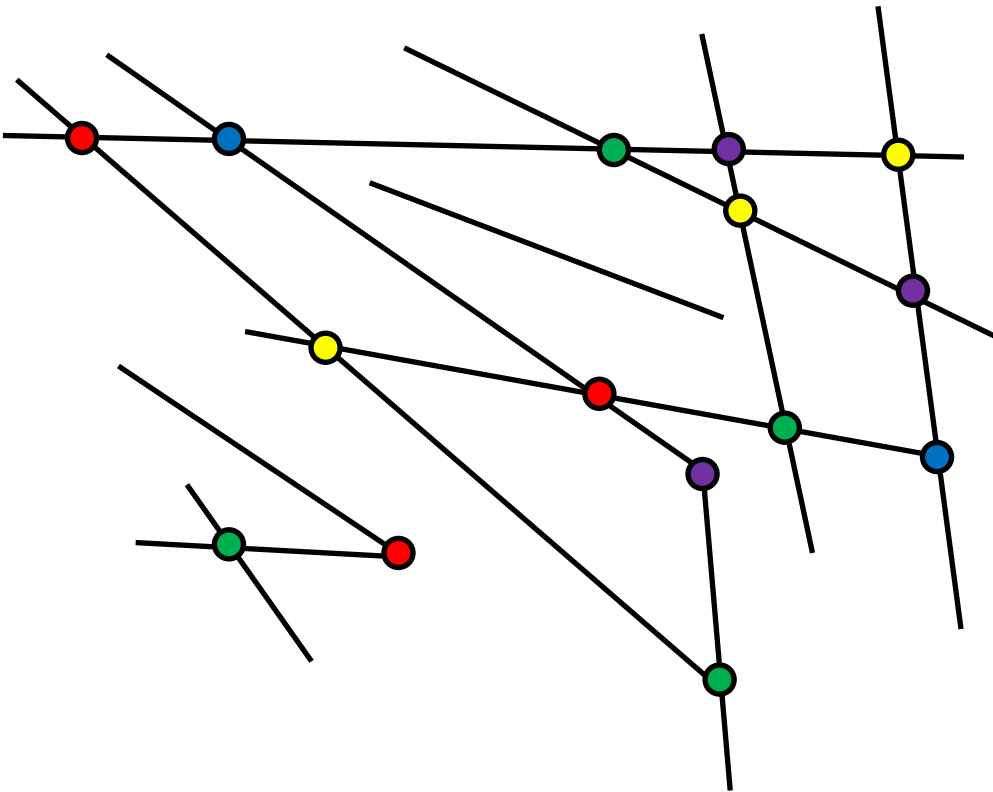
Corresponding graph

G_M



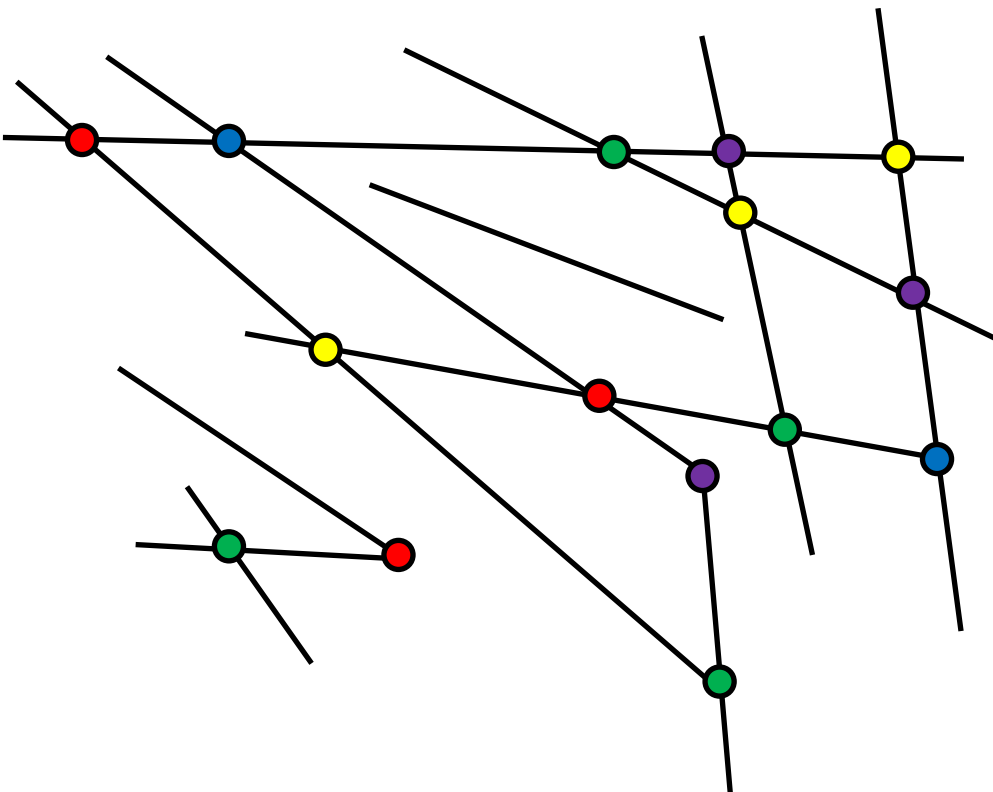
Coloring a set of segments

Coloring a set of segments means coloring the intersection points so that no segment has points with the same color.



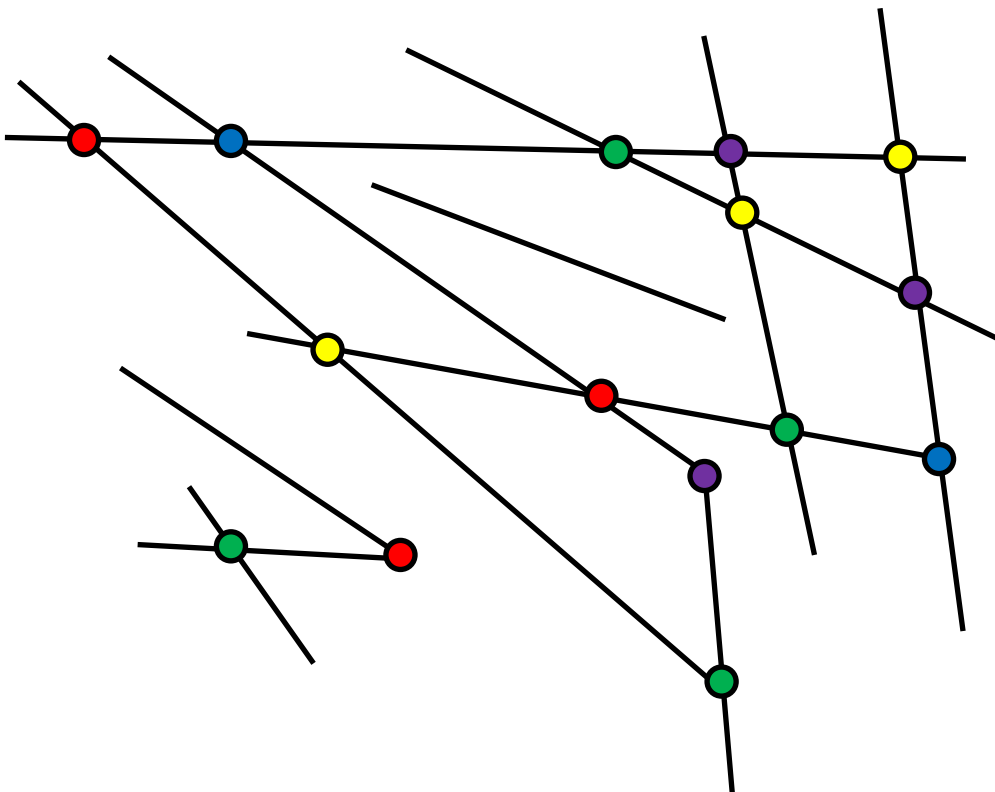
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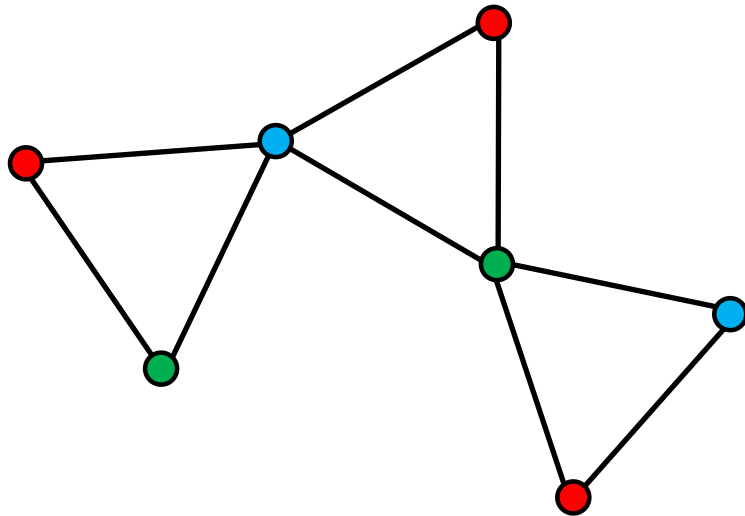
$$\chi(M) = 5$$

Erdős-Faber-Lovász Conjecture

Let G be a graph consisting of n copies of K_n , every pair of which has at most one vertex in common. Then, $\chi(G) = n$.

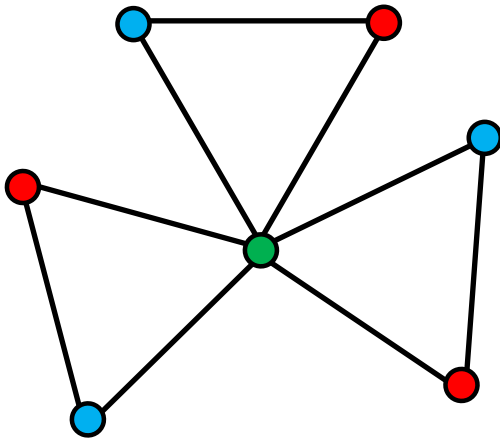
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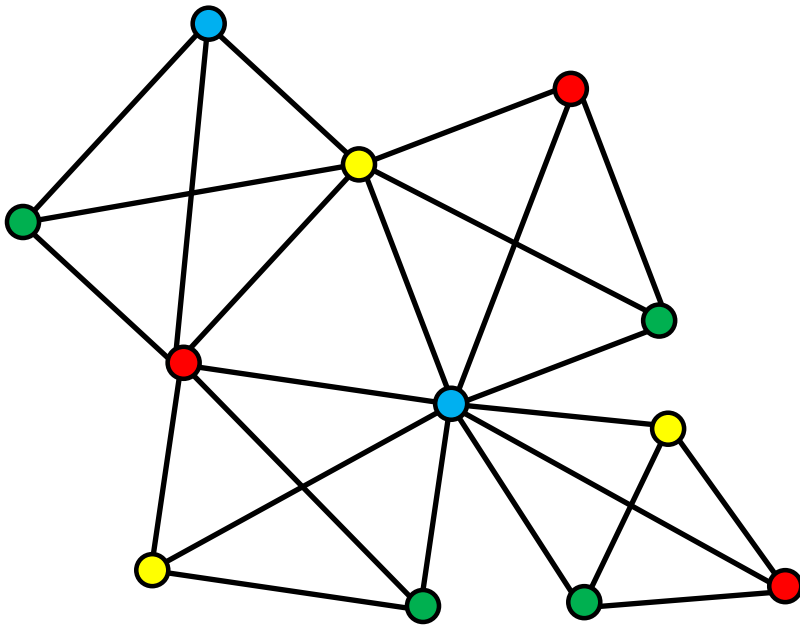
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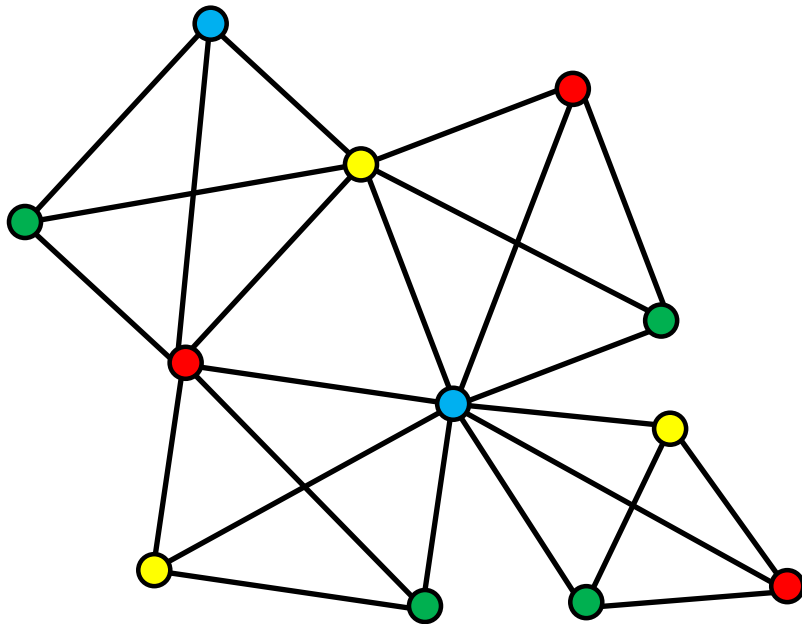
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Erdős-Faber-Lovász Conjecture

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$\chi(G)$ cannot be less than n , since at least n colors are needed for each copy of K_n

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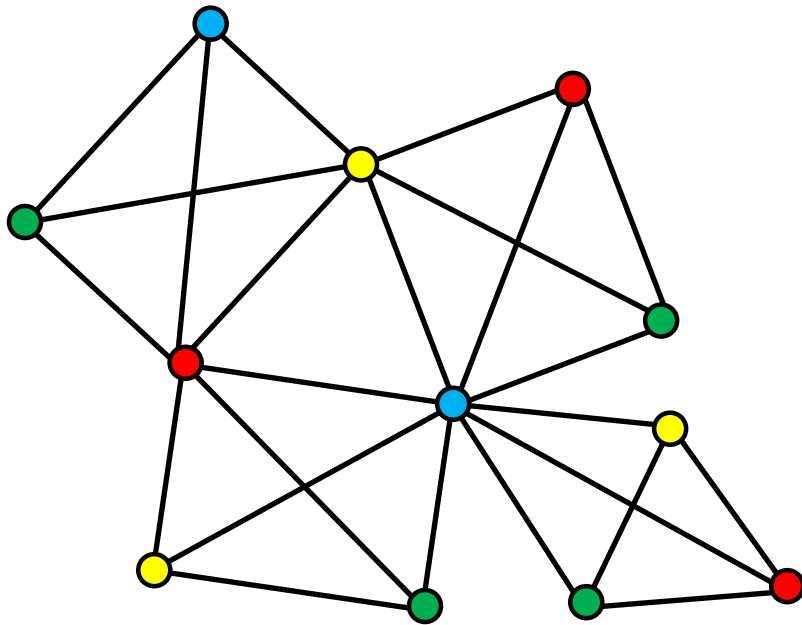
[Kang et al. 2021] announced a proof for large enough n

Erdős-Faber-Lovász Conjecture

Let M be a set of m curves, each pair of which has at most one point in common. Then, $\chi(M) \leq m$.

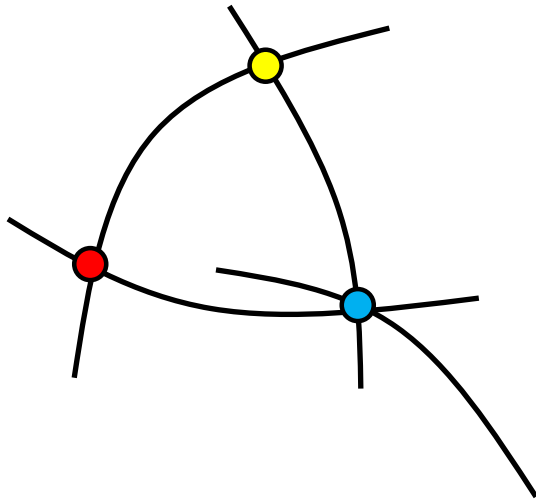
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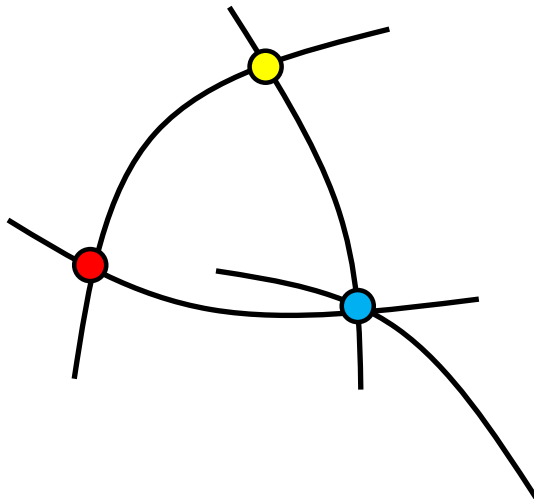
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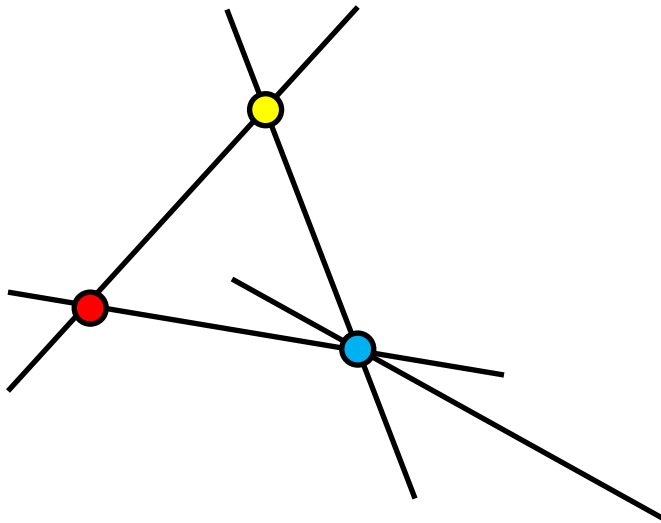
Let M be a set of m curves, each pair of which has at most one point in common. Then, $\chi(M) \leq m$.



$\chi(M)$ can be less than m .

Coloring a set of segments

Let M be a set of m segments. What's $\chi(M)$?




Bounds

$$w(M) \leq \chi(M) \leq |M|$$

Bounds

Need $w(M)$ colors
for the segment with
 $w(M)$ intersections

Follows from proof
of EFL Conjecture

$$w(M) \leq \chi(M) \leq |M|$$
The diagram consists of two blue arrows. The first arrow starts from the text 'Need w(M) colors for the segment with w(M) intersections' and points down to the first 'w(M)' in the inequality. The second arrow starts from the text 'Follows from proof of EFL Conjecture' and points down to the second 'w(M)' in the inequality.

Bounds

Need $w(M)$ colors
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Follows from proof
of EFL Conjecture

$$w(M) \leq \chi(M) \leq |M|$$

easy to compute



Complexity of $\chi(M)$

Computing $\chi(M)$ is NP-Complete, even if M is a set of segments where no five segments intersect in the same point.

Reduction using a special case of PLANAR GRAPH
COLORING

Bounds

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hard to compute

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Questions

Experimentally, $\chi(M)$ usually seems to be a lot closer to $w(M)$ than to $|M|$.

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 - Can $\chi(M) \geq w(M) + 3$?
- Characterizations are hard to derive
- Hard to find one

Questions

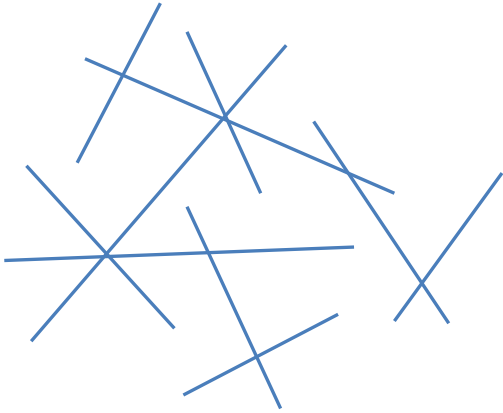
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What is $\chi(M)$ for M with special structure?

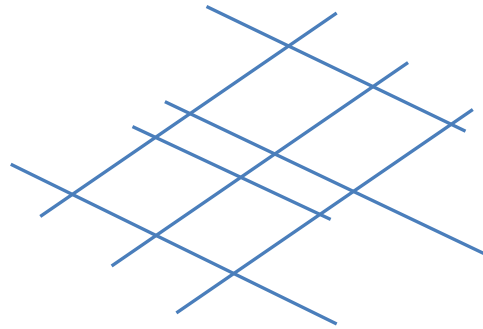
Easier to answer

Some characterizations



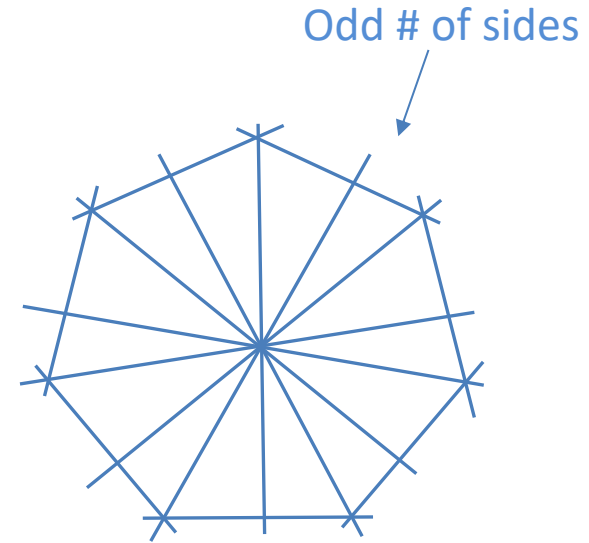
“circuit-free
arrangements”

$$\chi(M) = w(M)$$



“grid-like
arrangements”

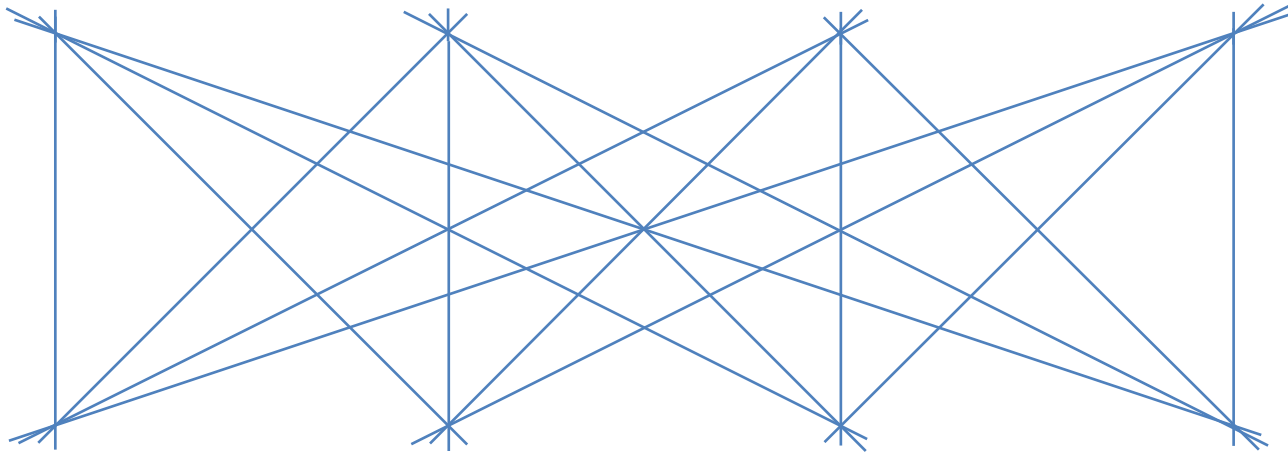
$$\chi(M) = w(M)$$



“bisected polygon
arrangements”

$$\chi(M) = w(M) + 2$$

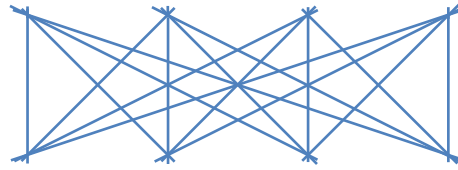
Some characterizations



“standard drawing of $K_{n,n}$ ”

$$\chi(M) = |\{(x, y): 1 \leq x, y \leq n, \gcd(x, y) = 1\}| + 2$$

Approach for



family

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Approach for family

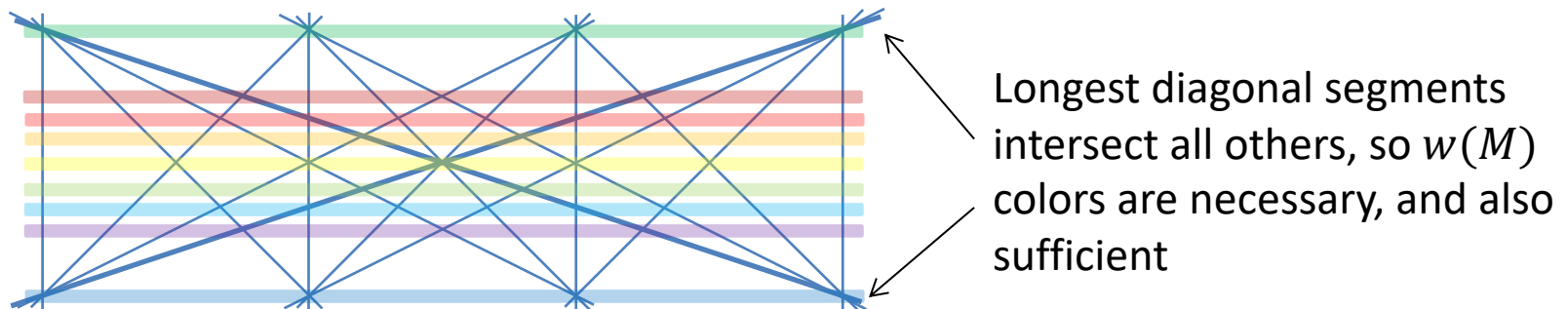
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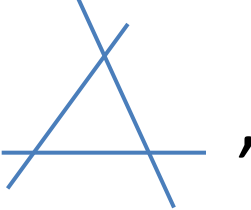
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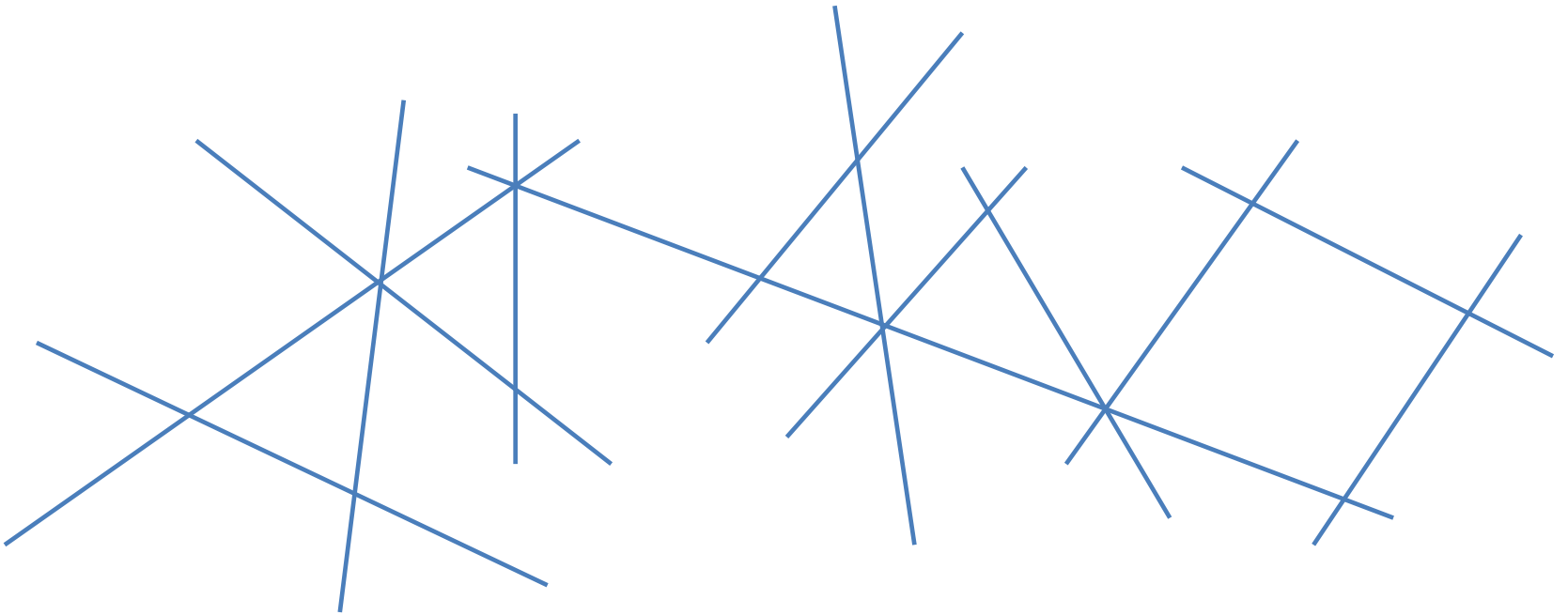
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Some characterizations

If M is a segment cactus different from , then $\chi(M) \leq |M| - 1$.



What about lines?

In a set M of lines, every pair of non-parallel lines intersect

$w(M) \leq \chi(M) \leq |M|$ still holds

What about lines?

The following are equivalent:

- If M is a set of m lines drawn in the plane, then M has a coloring with m colors.
- If M is a set of m segments drawn in the plane, then M has a coloring with m colors.

What about lines?

The following are equivalent:

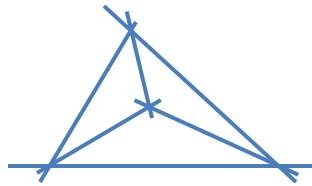
- If M is a set of m lines drawn in the plane, then M has a coloring with m colors.
- If M is a set of m segments drawn in the plane, then M has a coloring with m colors.

An analogous statement does not necessarily hold for subfamilies of segments/lines

Approach for family

Approach for family

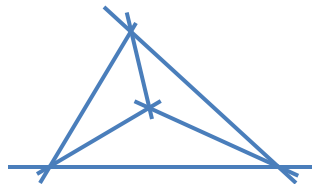
- Randomly noticed one small example



$$\chi(M) = w(M) + 2$$

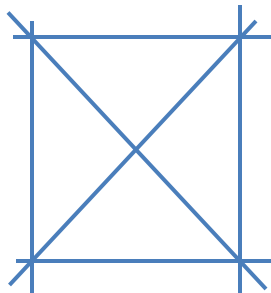
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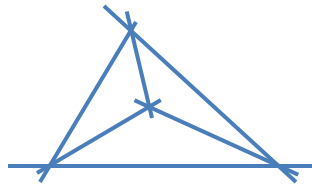
- Noticed another that was similar to it



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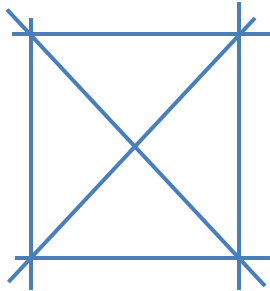
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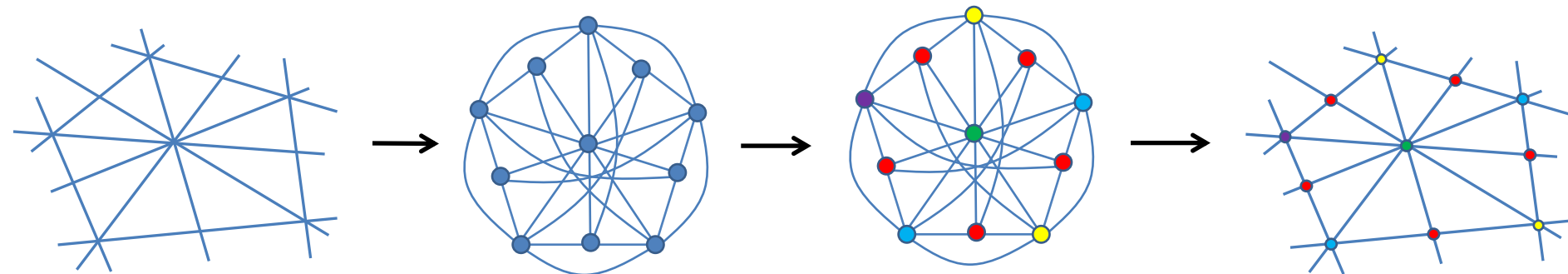


$$\chi(M) = w(M) + 2$$

- Tried to extend the pattern
- Harder to verify large examples by hand

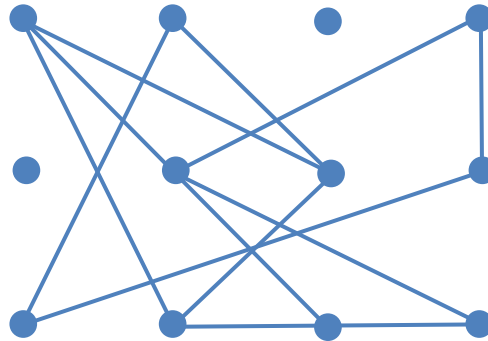
Approach for family

- Wrote integer program to help verification
- Used program to check more examples
- Saw $\chi(M) = w(M) + 2$ holds for all of them
- Analytically proved formula
 - Proof is based on converting the set of segments to a graph, finding a coloring that always works, and showing that no smaller coloring exists



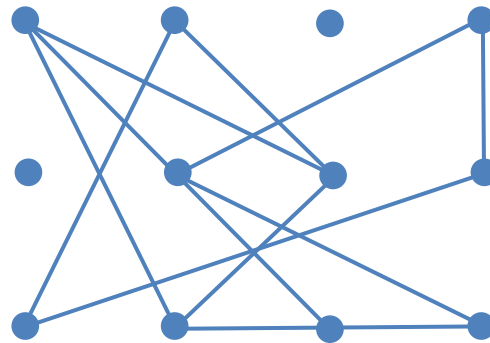
Future work

- Find a set with $\chi(M) \geq w(M) + 3$ or show this can't exist.
- Run more extensive computational experiments



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Thank you